An overview of ∞ -categories and higher algebra

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Why use ∞ -categories?

Some phenomena and propositions cannot be stated in full clarity without ∞ -categories.

Example

Ohromatic convergence and chromatic pullback:



Chromatic convergence and chromatic pullback should be described as homotopy limits of homotopy coherent diagrams $N(\mathbb{Z}_{\geq 0}^{op}) \to Sp$ and $\Lambda_2^2 \to Sp$ instead of homotopy diagrams $\mathbb{Z}_{\geq 0}^{op} \to h(Sp)$ or $\Lambda_2^2 \to h(Sp)$.

② Similarly, a Postnikov tower in the category $\mathcal S$ of spaces and its convergence.

Example (More)

- If C is a 1-category, then $Sp(C) \simeq \{*\}$ is trivial. The stabilization for 1-categories is meaningless. **Stable homotopy** is a higher categorical phenomenon.
- ② By ∞-categories we can define all kinds of moduli spaces, such as $CAlg(Sp) \times_{CAlg(hSp)} \{R\}$, the moduli space of \mathbb{E}_{∞} -structures on a given homotopy commutative ring spectrum R. The \mathbb{E}_{∞} -structures on a Lubin–Tate spectrum $E(n, \Gamma)$ is unique, meaning $CAlg(Sp) \times_{CAlg(hSp)} \{E(n, \Gamma)\}$ is a contractible Kan complex.
- **Observed Solution** Bousfield localization and connective cover of an \mathbb{E}_{∞} -ring are still \mathbb{E}_{∞} -rings. In the ∞ -categorical setting, this is automatic by the fact $L_E : Sp \rightleftharpoons Sp_E : i$ and $i : Sp_{\geq 0} \rightleftharpoons Sp : \tau_{\geq 0}$ are symmetric monoidal adjunctions, which induce adjunctions $CAlg(Sp) \rightleftharpoons CAlg(Sp_E)$ and $CAlg(Sp_{\geq 0}) \rightleftharpoons CAlg(Sp)$.
- Equivariant stable homotopy theory: there are numerous model categories characterizing it, but all of their underlying ∞-categories are equivalent to Fun(BG, Sp), which is both simple and intuitive.

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- G Equivariant stable homotopy theory: there are numerous model categories characterizing it, but all of their underlying ∞-categories are equivalent to Fun(BG, Sp), which is both simple and intuitive.

Motivation

The most significant motivation is to enrich the morphism set $Hom_{\mathcal{C}}(X, Y)$ in a category \mathcal{C} to a topological space $Map_{\mathcal{C}}(X, Y)$. Then we can have higher morphisms $\pi_n Map_{\mathcal{C}}(X, Y)$.

For example, when considering the category of spectra, we have $\pi_n Map_{\mathcal{C}}(X, Y) = [\Sigma^n X, Y] = Y^{-n}(X).$

So the most intuitive model for ∞ -category theory should be *sSet*-enriched (or *Top*-enriched) categories. However, all of these models are equivalent to Joyal's model. Indeed we have Quillen equivalences $(sSet)_{Joyal} \rightleftharpoons Cat_{sSet} \rightleftharpoons Cat_{Top}$.

But Joyal's model encodes information more concisely: the only data of a quasi-category is a simplicial set.

Underlying \mathcal{H} -enriched category

There are many different ways to extract mapping spaces $Map_{\mathcal{C}}(X, Y)$ from an ∞ -category \mathcal{C} .

But when we take their underlying $\mathcal{H} := Ho(sSet_{Kan})$ -enriched categories, all of them become the same, written as <u> $h\mathcal{C}$ </u>.

Remark

The processes $\mathcal{C} \mapsto \underline{h\mathcal{C}} \mapsto \underline{h\mathcal{C}}$ make it simpler to manage but meanwhile cause a loss of homotopy coherent information. How to extract useful and discard redundant information of homotopy coherence in specific circumstances is an "art" in ∞ -categories' world.

Preventing Russell's paradox

In order to consider the **category of all categories**, we need to add a set-theoretic axiom into ZFC, i.e. Grothendieck's Assumption:

∀ cardinal κ , there exists an inaccessible cardinal $\tau > \kappa$. (A good reference: Chap 1, 代数学方法 1, 李文威)

Methodology

By Grothendieck's Assumption,

1. When not involving **category of all categories**, technically we can treat all things as small. So all propositions not involving **category of all categories** will hold in any Grothendieck universe.

 When involving category of all categories, for example Cat_∞, we consider it as the ∞-category Cat_∞^τ of all τ-small categories for an inaccessible cardinal τ. Choose a bigger inaccessible τ₂ > τ, then technically we can treat Cat_∞^τ as a τ₂-small ∞-category in Cat_∞^{τ₂}.

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2. When involving **category of all categories**, for example Cat_{∞} , we consider it as the ∞ -category Cat_{∞}^{τ} of all τ -small categories for an inaccessible cardinal τ . Choose a bigger inaccessible $\tau_2 > \tau$, then technically we can treat Cat_{∞}^{τ} as a τ_2 -small ∞ -category in $Cat_{\infty}^{\tau_2}$.

Definition (Kan extension along a full subcategory)

Let $i : \mathcal{C}_0 \subset \mathcal{C}$ be a full subcategory, we say a functor $F : \mathcal{C} \to \mathcal{D}$ is a left Kan extension along i iff $\forall X \in \mathcal{C}$, $(\mathcal{C}_0 \times_{\mathcal{C}} \mathcal{C}_{/X})^{\triangleright} \to \mathcal{C} \xrightarrow{F} \mathcal{D}$ is a colimit diagram, i.e. $colim_{A \to X, A \in \mathcal{C}_0} F(A) \simeq F(X)$.

Theorem

The restriction $Fun^{LKan}(\mathcal{C}, \mathcal{D}) \xrightarrow{\sim} Fun^{\exists LKan}(\mathcal{C}_0, \mathcal{D})$ is a categorical equivalence.

Example

Let C be a small category and D be a category that admits small colimits, then (1) A functor $F : \mathcal{P}(C) \to D$ is a left Kan extension along the Yoneda embedding $i : C \to \mathcal{P}(C)$ iff F preserves small colimits. (2) For any $f \in Fun(C, D)$, there exists a left Kan extension $F : \mathcal{P}(C) \to D$ along (3) And hence we have $Fun^{colim}(\mathcal{P}(C), D) \to Fun(C, D)$ is an equivalence. (e.g.

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Definition

Let \mathbb{K} be a collection of simplicial sets. We say that an ∞ -category \mathcal{C} is \mathbb{K} -cocomplete if it admits K-diagram colimits, for each $K \in \mathbb{K}$.

We say that a functor of ∞ -categories $h: \mathcal{C} \to \widehat{\mathcal{C}}$ exhibits $\widehat{\mathcal{C}}$ as a K-cocompletion of \mathcal{C} if the ∞ -category $\widehat{\mathcal{C}}$ is K-cocomplete and for every K-cocomplete ∞ -category \mathcal{D} , precomposition with h induces an equivalence of ∞ -categories $\operatorname{Fun}^{\mathbb{K}}(\widehat{\mathcal{C}}, \mathcal{D}) \xrightarrow{\sim} \operatorname{Fun}(\mathcal{C}, \mathcal{D}).$

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Let \mathbb{K} be a (small) collection of simplicial sets, then for any (small) ∞ -category C, there exists a \mathbb{K} -completion $C \to P^{\mathbb{K}}(C)$. That gives an adjunction $\widehat{Cat}_{\infty} \rightleftharpoons \widehat{Cat}(\mathbb{K})_{\infty}$, e.g. $P^{small}(C) = Fun(C, S)$ and $P^{small}(*) = S$.

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Let \mathcal{D} be an ∞ -category.

Theorem (Pointedlization)

If \mathcal{D} admits final object, then there exists a pointedlization $\mathcal{D}_{*/} \to \mathcal{D}$ such that for any pointed ∞ -category \mathcal{C} the forgetful functor θ : Fun' $(\mathcal{C}, \mathcal{D}_*) \to \operatorname{Fun'}(\mathcal{C}, \mathcal{D})$ is an equivalence. That provides an adjunction $\operatorname{Cat}_{\infty}^{Final, pt} \rightleftharpoons \operatorname{Cat}_{\infty}^{Final}$.

Theorem (Stabilization)

If \mathcal{D} admits finite limits, then there exists a stabilization $Sp(\mathcal{D}) \to \mathcal{D}$ such that for any stable ∞ -category \mathcal{C} the forgetful functor $\theta : \operatorname{Fun}^{Flim}(\mathcal{C}, Sp(\mathcal{D})) \to \operatorname{Fun}^{Flim}(\mathcal{C}, \mathcal{D})$ is an equivalence. That provides an adjunction $Cat_{\infty}^{Flim,st} \rightleftharpoons Cat_{\infty}^{Flim}$.

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The category spectra Sp(P(*)) is the stabilization of the cocompletion of the trivial ∞ -category.

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Let $n \ge -2$, an object Z in an ∞ -category C is *n*-truncated if, for every object $Y \in C$, the space $Map_C(Y, Z)$ is *n*-truncated space.

Theorem (Truncation)

If C is a presentable ∞ -category, then there exists an *n*-truncation functor $C \to \tau_{\leq n} C$. Suppose that \mathcal{D} is a presentable that all objects are *n*-truncated, i.e. it's an (n + 1)-category. Then composition with $\tau_{\leq n}$ induces an equivalence $s : \operatorname{Fun}^{\mathrm{L}}(\tau_{\leq n} \mathcal{C}, \mathcal{D}) \to \operatorname{Fun}^{\mathrm{L}}(\mathcal{C}, \mathcal{D})$. That provides an adjunction $Pr^{L} \rightleftharpoons Pr^{L}_{\leq (n+1)}$.

Example

(1) An space X in S is n-truncated iff all $\pi_i X$ vanish when i > n. Particularly $S_{\leq 0} \simeq N(Set)$. (2) An n-truncated object Cat_{∞} is exactly an n-category. And all n-categories form an (n+1)-category $(Cat_{\infty})_{\leq n}$.

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Postnikov-type tower

Let \mathcal{C} be an ∞ -category and $I = \{\mathcal{C}_0 \subset \mathcal{C}_1 \subset \cdots \subset \mathcal{C}_n \subset \cdots \subset \mathcal{C}\}$ be an ascending sequence of reflective full subcategories of \mathcal{C} , where **reflective** means the inclusion functor $\mathcal{C}_i \hookrightarrow \mathcal{C}$ admits a left adjoint.

Example

- If taking $I = \{S_{\leq 0} \subset S_{\leq 1} \subset \cdots \subset S_{\leq n} \subset \cdots \subset S\}$, we recover to the classical case.
- ② If taking $I = \{L_0 Sp^{\omega}_{(p)} \subset L_1 Sp^{\omega}_{(p)} \subset \cdots \subset L_n Sp^{\omega}_{(p)} \subset \cdots \subset Sp^{\omega}_{(p)}\}$ where $Sp^{\omega}_{(p)}$ is the ∞-category of finite *p*-local spectra, we get chromatic convergence case.

Definition (Tower and pretower)

- An *I*-tower in C is a functor $N(\mathbb{Z}_{\geq 0}^{op})^{\triangleleft} \to C$, which we view as a diagram $X_{\infty} \to \cdots \to X_2 \to X_1 \to X_0$ satisfying that for each $n \geq 0$, the map $X_{\infty} \to X_n$ exhibits X_n as a C_n -reflection of X_{∞} .
- ② An *I*-pretower in *C* is a functor $N(\mathbb{Z}_{\geq 0}^{op}) \to C$: ... → $X_2 \to X_1 \to X_0$ which exhibits each X_n as a C_n -reflection of X_{n+1} .

Postnikov-type convergence

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Definition

We let $\operatorname{Post}_{I}^{+}(\mathcal{C})$ denote the ∞ -category of *I*-towers, and $\operatorname{Post}_{I}(\mathcal{C})$ the ∞ -category of *I*-pretowers. We have an evident forgetful functor $\phi : \mathcal{C} \xleftarrow{} \operatorname{Post}_{I}^{+}(\mathcal{C}) \to \operatorname{Post}_{I}(\mathcal{C})$. We will say that \mathcal{C} is **Postnikov** *I*-complete if ϕ is an equivalence of ∞ -categories.

Theorem (Postnikov-type convergence)

Suppose that any *I*-pretower in *C* has a limit. Then *C* is Postnikov *I*-complete **if and only if**, for every diagram $X : \mathbb{N}(\mathbb{Z}_{\geq 0}^{op})^{\triangleleft} \to C$ the following conditions are equivalent: (1) The diagram X is an *I*-tower. (2) The diagram X is a limit in *C*, and the restriction of X to $\mathbb{N}(\mathbb{Z}_{\geq 0}^{op})$ is an *I*-pretower.

Definition (Reformulation of ordinary commutative monoids)

A (3-)commutative monoid in an ordinary category \mathcal{C} which admits finite products is a functor $M: (Fin_*)_{\leq 3} \to \mathcal{C}$ such that the canonical maps $M(\rho_i): M(\langle n \rangle) \to M(\langle 1 \rangle)$ exhibit $M(\langle n \rangle) \simeq \prod_{1 \leq i \leq n} M(\langle 1 \rangle)$ in \mathcal{C} for all $0 \leq n \leq 3$.



Definition (∞ -commutative monoid)

Let \mathcal{C} be an ∞ -category with finite products. We define an ∞ -commutative monoid in \mathcal{C} as a functor $M: N(Fin_*) \to \mathcal{C}$ such that the canonical maps $M(\rho_i): M(\langle n \rangle) \to M(\langle 1 \rangle)$ exhibit $M(\langle n \rangle) \simeq \prod_{1 \le i \le n} M(\langle 1 \rangle)$ in \mathcal{C} for all $n \ge 0$.

Proposition (Barkan 2022)

Let \mathcal{C} be a complete *n*-category. Then $CMon^{\infty}(\mathcal{C}) \xrightarrow{\sim} CMon^{n+2}(\mathcal{C})$ is categorically equivalent.

Definition

A symmetric monoidal ∞ -category is an (∞ -)commutative monoid in Cat_{∞} .

Corollary

Particularly, if a symmetric monoidal ∞ -category C is a 1-category, then it is an ∞ -commutative monoid in $(Cat_{\infty})_{\leq 1}$, which is a 2-category and written as $Cat_{\leq 1}$. So we have $CMon^{\infty}(Cat_{\leq 1}) \xrightarrow{\sim} CMon^4(Cat_{\leq 1})$.

It can be checked that the 4-commutativity in $Cat_{\leq 1}$ exactly corresponds with ordinary coherent conditions of a symmetric monoidal category.

Lurie's definition

By the (un)straightening equivalence $Fun(N(Fin_*), Cat_{\infty}) \simeq CoCart_{/N(Fin_*)}$, we get the following equivalent definition by Lurie.

Definition

A symmetric monoidal ∞ -category is a coCartesian fibration of simplicial sets $p: \mathcal{C}^{\otimes} \to N(Fin_*)$ with the following property: For each $n \geq 0$, the maps $\{\rho^i : \langle n \rangle \to \langle 1 \rangle\}_{1 \leq i \leq n}$ induce functors $\rho^i_! : \mathcal{C}^{\otimes}_{\langle n \rangle} \to \mathcal{C}^{\otimes}_{\langle 1 \rangle}$ which determine an equivalence $\mathcal{C}^{\otimes}_{\langle n \rangle} \simeq (\mathcal{C}^{\otimes}_{\langle 1 \rangle})^n$. We define $\mathcal{C}^{\otimes}_{\langle 1 \rangle}$ as its underlying ∞ -category.

This definition has technical advantages for general ∞ -operads.

Tensor product of ∞ -categories

Let \mathbb{K} be the collection of all small simplicial sets.

Definition

Given 2 cocomplete ∞ -categories C and D, we define the tensor product as a functor $C \times D \to C \otimes D$ such that for any cocomplete E, we have $Fun^{\mathbb{K}}(C \otimes D, E) \xrightarrow{\sim} Fun^{\mathbb{K}\boxtimes\mathbb{K}}(C \times D, E)$. Such tensor product always exists because the natural functor $C \times D \to \mathcal{P}_{\mathbb{K}\boxtimes\mathbb{K}}^{\mathbb{K}}(C \times D)$ satisfies that.

Theorem

The above gives a symmetric monoidal structure $\widehat{Cat}_{\infty}(\mathbb{K})^{\otimes} \to N_{*}(Fin_{*})$ and makes the cocompletion funcor a symmetric monoidal adjunction $\widehat{Cat}_{\infty}^{\otimes} \rightleftharpoons \widehat{Cat}_{\infty}(\mathbb{K})^{\otimes}$. So $\mathcal{S} = \mathcal{P}(*)$ is the unit in $\widehat{Cat}_{\infty}(\mathbb{K})^{\otimes}$.

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Remark

By (un)straightening equivalence, $CAl(\widehat{Cat}_{\infty}(\mathbb{K})) \subset CAl(\widehat{Cat}_{\infty})$ is the subcategory whose objects are symmetric monoidal ∞ -categories such that $-\otimes -$ preserves colimits separately in each variable (called **cocomplete symmetric monoidal** categories), and whose morphisms are **colimit-preserving** symmetric monoidal functors.

Corollary

The symmetric monoidal adjunction induces an adjunction between algebras $F: CAl(\widehat{Cat}_{\infty}) \rightleftharpoons CAl(\widehat{Cat}_{\infty}(\mathbb{K})).$

Corollary

(1) The $S = \mathcal{P}(*)$ is the unit in $\widehat{Cat}_{\infty}(\mathbb{K})^{\otimes}$, which means it is initial object in $CAl(\widehat{Cat}_{\infty}(\mathbb{K}))$ and hence S admits a cocomplete symmetric monoidal structure S. (2) So for any cocomplete symmetric monoidal ∞ -category, there exists essentially unique colimit-preserving symmetric monoidal functor $S^{\otimes} \to C^{\otimes}$.

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Proposition (Localization)

Let C be an ∞ -category and let $L : C \to C$ be a functor with essential image $LC \subseteq C$. The following conditions are equivalent: (1) There exists a functor $f : C \to D$ with a fully faithful right adjoint $g : D \to C$ and

an equivalence between $g \circ f$ and L.

(2) When regarded as a functor from C to LC, L is a left adjoint of the inclusion $LC \subseteq C$.

(3) There exists a natural transformation from $id_{\mathcal{C}} \to L$ such that, $L \circ id_{\mathcal{C}} \to L \circ L$ and $id_{\mathcal{C}} \circ L \to L \circ L$ are equivalences in $\operatorname{Fun}(\mathcal{C}, \mathcal{C})$, i.e. an idempotent object in $\operatorname{Fun}(\mathcal{C}, \mathcal{C})$.

Proposition

The full subcat $Pr^{L} \subset \widehat{Cat_{\infty}}(\mathbb{K})$ is closed under tensor product and hence inherits a symmetric monoidal structure Pr_{L}^{\otimes} .

Proposition (Symmetric monoidal colocalization)

Let $C^{\otimes} \to N(Fin_*)$ be a symmetric monoidal ∞ -category. Let $\mathcal{D} \subseteq C$ be a full subcategory which is stable under equivalence. Suppose that the functor $- \otimes - : \mathcal{C} \times \mathcal{C} \to \mathcal{C}$ carries $\mathcal{D} \times \mathcal{D}$ into \mathcal{D} (meaning \mathcal{D} is **closed under tensor products**). Then the following hold.

- **1** The restricted map $\mathcal{D}^{\otimes} \to N(Fin_*)$ is a symmetric monoidal ∞ -category.
- **2** The inclusion $\mathcal{D}^{\otimes} \subseteq \mathcal{C}^{\otimes}$ is a symmetric monoidal functor.
- Suppose that the inclusion D ⊆ C admits a right adjoint L (so that D is a colocalization of C). Then there exists a lax-symmetric-monoidal right adjunction L[⊗] : C[⊗] → D[⊗].

Formally speaking, L^{\otimes} is a right adjunction in the strict 2-category $h_2(Op_{O\otimes})$.

Corollary

Under assumptions of (3) above, a symmetric monoidal colocalization can induce a colocalization on algebras $CAlg(\mathcal{D}) \rightleftharpoons CAlg(\mathcal{C})$.

Corollary (*t*-structure and symmetric monoidal structure)

Let $p: \mathcal{C}^{\otimes} \to N(Fin_*)$ be a symmetric monoidal ∞ -category. Assume that the underlying ∞ -category \mathcal{C} is stable and that $-\otimes$ – is exact in each variable. We will say that a *t*-structure ($\mathcal{C}_{\geq 0}, \mathcal{C}_{\leq 0}$) is **compatible** with the symmetric monoidal structure if the functor \otimes carries $\mathcal{C}_{\geq 0} \times \mathcal{C}_{\geq 0}$ into $\mathcal{C}_{\geq 0}$. Then the induced map $\mathcal{C}_{\geq 0}^{\otimes} \to N(Fin_*)$ is again a symmetric monoidal ∞ -category,

and $\mathcal{C}_{\geq 0}^{\otimes} \xleftarrow{i}{\tau_{\geq 0}} \mathcal{C}^{\otimes}$ is a symmetric monoidal colocalization. Thus it further induces a colocalization $CAlg(\mathcal{C}_{\geq 0}) \xleftarrow{i}{\tau_{\geq 0}} CAlg(\mathcal{C})$.

Example (Connective cover of an \mathbb{E}_{∞} -ring)

When $\mathcal{C} = Sp$ we have $CAlg(Sp_{\geq 0}) \xleftarrow[\tau_{\geq 0}]{} CAlg(Sp)$, which means that the connective cover of an \mathbb{E}_{∞} -ring naturally inherits an \mathbb{E}_{∞} -structure.

Proposition (Symmetric monoidal localization)

Let $C^{\otimes} \to N(Fin_*)$ be a symmetric monoidal ∞ -category. Let $\mathcal{D} \subseteq C$ be a full subcategory. Suppose that $\mathcal{D} \subset C$ is a reflective subcategory (with a left adjoint $L: C \to \mathcal{D}$). If for every pair g_1, g_2 of L-equivalences in C, the morphism $g_1 \otimes g_2$ in C is also an L-equivalence (meaning L-equivalences are closed under tensor products), then we have the following.

- **1** The restricted map $\mathcal{D}^{\otimes} \to N(Fin_*)$ is lax-symmetric-monoidal.
- **2** The inclusion $\mathcal{D}^{\otimes} \subseteq \mathcal{C}^{\otimes}$ is a symmetric monoidal functor.

③ There exists a symmetric monoidal left adjoint $L^{\otimes} : \mathcal{C}^{\otimes} \to \mathcal{D}^{\otimes}$.

Corollary

A symmetric monoidal localization can induce a localization on algebras $CAlg(\mathcal{C}) \rightleftharpoons CAlg(\mathcal{D}).$

Bousfield localization

Let \mathcal{C}^{\otimes} be a presentably symmetric monoidal ∞ -category, i.e., an object in $CAlq(Pr^L) \hookrightarrow CAlq(Cat_{\infty}).$

Theorem (Bousfield localization)

Let $E \in \mathcal{C}$ be an object. Then $W_E = \{X \to Y | X \otimes E \xrightarrow{\sim} Y \otimes E\} \subset Fun(\Delta^1, \mathcal{C})$ is a small-generated strongly saturated collection of morphisms, which means that there exists an accessible localization functor $L_E: \mathcal{C} \to \mathcal{C}$.

Furthermore, Bousfield localization is compatible with its symmetric monoidal

structure, meaning it forms a symmetric monoidal localization $\mathcal{C}^{\otimes} \xleftarrow{L_E^{\otimes}}{\mathcal{C}_E^{\otimes}}$.

Example (Bousfield localization of an \mathbb{E}_{∞} -ring)

When $\mathcal{C} = Sp$ we have $CAlg(\mathcal{C}) \xleftarrow{CAlg(\mathcal{L}_E)} CAlg(\mathcal{C}_E)$. This means that Bousfield localization of an \mathbb{E}_{∞} -ring naturally inherits an \mathbb{E}_{∞} -structure.

Idempotent object

Let \mathcal{C} be a symmetric monoidal ∞ -category.

Definition (idempotent object)

Let $e: 1_C \to E$ be a morphism in C. We say e is idempotent iff $1_C \otimes X \to X \otimes X$ is equivalent. (e.g. $\mathbb{Z} \to \mathbb{Z}[1/p]$ in Ab)

Theorem (Bousfield localization with respect to an idempotent object)

Let $e: 1_C \to E$ be a morphism in \mathcal{C} , then

(1) The *e* is an idempotent object of *C* iff the transformation $\alpha : id_{\mathcal{C}} \to l_E$ exhibits l_E as a localization functor on *C*, where $l_E : \mathcal{C} \to \mathcal{C}$ is given by the tensor product with *E*. (2) If *e* is idempotent, then l_E is exactly the Bousfield localization with respect to *E*, which has the following properties:

(a) The l_E is compatible with \otimes , so induces a symmetric monoidal localization

$$\mathcal{C}^{\otimes} \xleftarrow{L_E^{\otimes}}{i^{\otimes}} \mathcal{C}_E^{\otimes}$$
 ,

(b) The inclusion i^{\otimes} is also symmetric monoidal, meaning C_E is closed under tensor products.

Definition

Let \mathcal{C} be a symmetric monoidal ∞ -category. We will say that a commutative algebra object $A \in \operatorname{CAlg}(\mathcal{C})$ is idempotent if unit map $e : \mathbf{1} \to A$ is idempotent.

Theorem

Let C be a symmetric monoidal ∞ -category with unit object 1, which we regard as a trivial algebra object of C. Then the functor

 $\theta : \operatorname{CAlg}^{idem}(\mathcal{C}) \subseteq \operatorname{CAlg}(\mathcal{C}) \simeq \operatorname{CAlg}(\mathcal{C})_{1/} \to \mathcal{C}_{1/}$

is fully faithful, and its essential image are idempotent objects in \mathcal{C} , which gives an equivalence $\operatorname{CAlg}^{idem}(\mathcal{C}) \xrightarrow{\sim} (\mathcal{C}_{1/})^{idem}$. Furthermore, any mapping space in $(\mathcal{C}_{1/})^{idem}$ is either empty or contractible, i.e. $(\mathcal{C}_{1/})^{idem}$ is equivalent to a partial-order set N(I).

Interesting applications after the internalization

Proposition

The full subcat $Pr^{L} \subset \widehat{Cat}_{\infty}(\mathbb{K})$ is closed under tensor products (S is also the unit in Pr^{L}) and hence inherits a symmetric monoidal structure. In fact, for any $\mathcal{C}, \mathcal{D} \in Pr^{L}$, we have a natural equivalence $\mathcal{C} \otimes D \simeq RFun(\mathcal{C}^{op}, \mathcal{D})$.

Theorem (Unique symmetric monoidal structure)

The following 4 colimit-preserving functors $S \xrightarrow{\tau \leq n} \tau_{\leq n} S$, $S \xrightarrow{(-)_+} S_*$, $S \xrightarrow{\Sigma_+^{\infty}} Sp$, and $S \xrightarrow{*\mapsto \mathbb{Z}} N(Ab)$ are idempotent objects in Pr^L . Hence by $\operatorname{CAlg}(Pr^L)^{idem} \xrightarrow{\sim} (Pr^L_{S/})^{idem}$ we conclude that S resp. $S_{\leq n}$, S_* , Sp, N(Ab) only admits a unique cocomplete symmetric monoidal structure with the unit * resp. *, S^0 , $\Sigma^{\infty}S^0$, \mathbb{Z} .

Interesting applications after the internalization

By Bousfield localization with respect to idempotent objects, we have:

Corollary

The following 4 full subcategories of Pr^L are closed under tensor products. (a) $Pr_{\leq n+1}^L$: the ∞ -category of presentable (n + 1)-categories; (b) Pr_*^L : the ∞ -category of presentable pointed ∞ -categories; (c) Pr_{st}^L : the ∞ -category of presentable stable ∞ -categories, known as tensor-triangulated ∞ -categories or tt- ∞ -categories; (d) Pr_{1ad}^L : the ∞ -category of presentable additive 1-categories.

Corollary

The localization functors $Pr^L \xrightarrow{-\otimes \tau_{\leq n} S} Pr^L_{\leq n+1}$, $Pr^L \xrightarrow{-\otimes S_*} Pr^L_*$, $Pr^L \xrightarrow{-\otimes Sp} Pr^L_{st}$, and $Pr^L \xrightarrow{-\otimes N(Ab)} Pr^L_{1\text{-}ad}$ correspond with the *n*-truncation, copointedlization, costabilization, and 1-coadditivalization of presentable ∞ -categories respectively.

My specific interests

Use higher algebra and spectral algebraic geometry (SAG) to explore various intersections between homotopy theory and algebraic geometry.

Example

- A recent good example is the Chromatic Nullstellensatz by Burklund, Schlank, and Yuan. They proved that "algebraically closed" \mathbb{E}_{∞} -rings in $CAlg(Sp_{T(n)})$ are exactly those Lubin–Tate spectra E(L) with L an algebraically closed field. And for any non-zero T(n)-local \mathbb{E}_{∞} -ring R, there exists a geometric point $R \to E(L)$.
- ② For another beautiful example, the Devinatz–Hopkins theorem $L_{K(n)}S \simeq E_n^{hG_n}$ can be interpreted as QCoh(Spf(E_n)/G_n) ≃ Sp_{K(n)} in (formal) SAG.
- In the framework of SAG, we can study spectral moduli problems: given an algebro-geomtric stack \mathcal{M}_0 , can we give an \mathbb{E}_{∞} -realization \mathcal{M} making $\pi_0 \mathcal{M} = \mathcal{M}_0$?

It is true when $\mathcal{M}_0 = \mathcal{M}_{ell}$ for the moduli stack of elliptic curves and when $\mathcal{M}_0 = \mathcal{X}_{K^p}$ for some of the Shimura stacks. Then taking the global sections of \mathbb{E}_{∞} -stacks respectively, we get TMF and TAF, which are intriguing \mathbb{E}_{∞} -rings.

Orientation theory from Thom spectra

- Quillen discovered a deep connection between (homotopy) complex orientation set $\operatorname{Or}_h(MU, E) := \operatorname{Hom}_{CAlg(hSp)}(MU, E)$ and formal group laws over E_* , which became the cornerstone of chromatic homotopy theory.
- After that, Ando-Hopkins-Strickland discovered a correspondence
 Or_h(MU⟨6⟩, E) → C³(P_E; I(0)) between MU⟨6⟩-orientations and cubical
 structures. By uniqueness of cubical structures on any line bundle of any abelian
 variety, we can endow a unique MU⟨6⟩-orientation to any elliptic cohomology
 theory.

-enhancement of orientations

When comes to \mathbb{E}_{∞} -orientation space $\operatorname{Or}_{\mathbb{E}_{\infty}}(Mf, R) := \operatorname{Map}_{CAlg(Sp)}(Mf, R)$, combining the Thom adjunction $\operatorname{Mon}_{\mathbb{E}_{\infty}}(\mathcal{S})_{/Pic(Sp)} \xleftarrow{M(-)} CAlg(Sp)$ and the infinite loop space machine $\operatorname{Mon}_{\mathbb{E}_{\infty}}^{gp}(\mathcal{S}) \simeq Sp_{\geq 0}$ we can produce many interesting results.



By this adjunction we can get the following theorem.

Theorem (Ando–Blumberg–Gepner–Hopkins–Rezk)

Let Mf be the Thom \mathbb{E}_{∞} -spectrum induced by a map $f: X \to \operatorname{pic}(Sp)$ in $Sp_{\geq 0}$ and let R be an \mathbb{E}_{∞} -ring. Then $\operatorname{Or}_{\mathbb{E}_{\infty}}(Mf, R)$ is a torsor over the H-space $\operatorname{Map}_{Sp}(X, gl_1(R))$, meaning $\operatorname{Or}_{\mathbb{E}_{\infty}}(Mf, R)$ is either empty or homotopy equivalent to $\operatorname{Map}_{Sp}(X, gl_1(R))$.

Example

Particularly, combining with the Chromatic Nullstellensatz and some further calculations, we can deduce that for any $\operatorname{height} = n > 0$, the $\operatorname{Or}_{\mathbb{E}_{\infty}}(MUP, E(\overline{\mathbb{F}}_p))$ is non-empty and hence homotopy equivalent to $\operatorname{Map}_{Sp}(ku, gl_1(E(\overline{\mathbb{F}}_p)))$.